

Characterizing Double Categories of Relations

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Axiomatization of Relations

A **relation** in a category is a monomorphism $R \rightarrow A \times B$.

\mathcal{C} = regular category

1. $\mathcal{C} \rightsquigarrow \text{Rel}(\mathcal{C})$ ordinary category/locally discrete 2-category
2. $\mathcal{C} \rightsquigarrow \mathfrak{R}el(\mathcal{C})$ a bicategory
3. $\mathcal{C} \rightsquigarrow \mathbb{R}el(\mathcal{C})$ a double category.

Question

- An **allegory** [FS90] is a locally-ordered 2-category equipped with an anti-involution $(-)^{\circ}$ satisfying a modularity law; every tabular allegory is equivalent to a category $\text{Rel}(\mathcal{C})$.
- A **bicategory of relations** [CW87] is a cartesian bicategory in which every object is discrete. Every functionally complete bicategory of relations is biequivalent to one of the form $\mathfrak{Rel}(\mathcal{C})$.
- Which double categories \mathbb{D} are equivalent to those of the form $\mathbb{R}\text{el}(\mathcal{C})$?

Double Categories

- A **double category** is a pseudo-category object in Cat .
- A double category \mathbb{D} is an **equipment** [Shu08] if every arrow f has a **companion** and a **conjoint**: proarrows f_{\downarrow} and f^* and cells

The diagram shows four commutative squares, each with a central downward-pointing arrow \Downarrow .
1. The first square has top horizontal arrow y and bottom horizontal arrow f_{\downarrow} . The left vertical arrow is 1 and the right vertical arrow is f .
2. The second square has top horizontal arrow f_{\downarrow} and bottom horizontal arrow y . The left vertical arrow is f and the right vertical arrow is 1 .
3. The third square has top horizontal arrow y and bottom horizontal arrow f^* . The left vertical arrow is f and the right vertical arrow is 1 .
4. The fourth square has top horizontal arrow f^* and bottom horizontal arrow y . The left vertical arrow is 1 and the right vertical arrow is f .

satisfying some equations [GP04].

- \mathbb{D} has **tabulators** if $y: \mathbb{D}_0 \rightarrow \mathbb{D}_1$ has a right adjoint.
- [Ale18] A double category \mathbb{D} is **cartesian** if $\Delta: \mathbb{D} \rightarrow \mathbb{D} \times \mathbb{D}$ and $\mathbb{D} \rightarrow 1$ have right adjoints in Dbl .
- Cartesian equipments $\text{Span}(\mathcal{C})$ characterized in [Ale18]
- Goal: characterize cartesian equipments of the form $\text{Rel}(\mathcal{C})$.

An **oplax/lax adjunction** is a conjoint pair in the strict double category of double categories with oplax and lax functors.

Theorem (§5 of [Nie12])

Let \mathbb{D} denote a double category with pullbacks. The following are equivalent:

1. *The identity functor $1: \mathbb{D}_0 \rightarrow \mathbb{D}_0$ extends to an oplax/lax adjunction $F: \text{Span}(\mathbb{D}_0) \rightleftarrows \mathbb{D}: G$.*
2. *\mathbb{D} is an equipment with tabulators.*

The oplax/lax adjunction is a strong (both functors are pseudo!) equivalence under some further completeness conditions:

Theorem (§5 of [Ale18])

For a double category \mathbb{D} the following are equivalent:

- 1. \mathbb{D} is equivalent to $\text{Span}(\mathcal{C})$ for finitely-complete \mathcal{C} .*
- 2. \mathbb{D} is a unit-pure cartesian equipment admitting certain Eilenberg-Moore objects.*
- 3. \mathbb{D}_0 has pullbacks satisfying a Beck-Chevalley condition and the canonical functor $\text{Span}(\mathbb{D}_0) \rightarrow \mathbb{D}$ is an equivalence of double categories.*

Relations as a Double Category

- Let $\mathcal{F} = (\mathcal{E}, \mathcal{M})$ be a proper and stable factorization system on a finitely complete category \mathcal{C} .
- The double category $\mathbb{R}el(\mathcal{C}; \mathcal{F})$ is a cartesian equipment.
- “Local products” [Ale18] in $\mathbb{R}el(\mathcal{C}; \mathcal{F})$ satisfy a Frobenius law:

$$R \wedge f^* \otimes S \cong f^* \otimes (f_! \otimes R \wedge S)$$

- $\mathbb{R}el(\mathcal{C}; \mathcal{F})$ is “unit-pure” i.e. $y: \mathbb{D}_0 \rightarrow \mathbb{D}_1$ is fully faithful.
- $\mathbb{R}el(\mathcal{C}; \mathcal{F})$ has tabulators.
- Eilenberg-Moore implies tabulators for spans. But tabulators are basic for allegories and bicategories of relations.

Tabulators

- An allegory is **tabular** if every arrow R has a tabulator: a pair of arrows f and g such that
 1. $gf^\circ = R$ “tabulators are strong”
 2. $f^\circ f \wedge g^\circ g = 1$ “tabulators are monic.”
- A double category \mathbb{D} has **tabulators** if $y: \mathbb{D}_0 \rightarrow \mathbb{D}_1$ has a right adjoint $\top: \mathbb{D}_1 \rightarrow \mathbb{D}_0$ in Dbl . The **tabulator** of $m: A \rightarrow B$ is the object $\top m$ together with a counit cell $\top m \Rightarrow m$.
- tabulators in $\text{Rel}(\mathcal{C}; \mathcal{F})$ satisfy:
 1. $\langle l, r \rangle: \top m \rightarrow A \times B$ is in \mathcal{M} and $l^* \otimes l_! \wedge r^* \otimes r_! \cong y$ holds (tabulators are monic)
 2. and $m \cong l^* \otimes r_!$ holds (strong).

Niefield's Theorem for Relations

Theorem

For a double category \mathbb{D} with a stable and proper factorization system $\mathcal{F} = (\mathcal{E}, \mathcal{M})$ on \mathbb{D}_0 , the identity functor $\mathbb{D}_0 \rightarrow \mathbb{D}_0$ extends to a normalized oplax/lax adjunction $F: \mathbb{R}el(\mathbb{D}_0; \mathcal{F}) \rightleftarrows \mathbb{D}: G$ if, and only if,

- 1. \mathbb{D} is a unit-pure equipment;*
- 2. has \mathcal{M} -monic tabulators;*
- 3. the unit cell y_e is an extension for each $e \in \mathcal{E}$.*

A First Pass on the Characterization

Theorem

The identity functor $1: \mathbb{D}_0 \rightarrow \mathbb{D}_0$ extends to an adjoint equivalence of pseudo-functors

$$F: \mathbb{R}el(\mathbb{D}_0; \mathcal{F}) \rightleftarrows \mathbb{D}: G$$

if, and only if,

1. \mathbb{D} is a unit-pure equipment;
2. y_e is an extension cell for each cover e ;
3. \mathbb{D} has strong, \mathcal{M} -monic and externally functorial tabulators;
4. every \mathcal{F} -relation $R \rightarrow A \times B$ is a tabulator of its canonical extension.

Covers and Inclusions

Definition (Cf. [Sch15])

The **kernel** of a morphism $f: A \rightarrow B$ is the restriction ρ of the unit on B along f . Dually, the **cokernel** of f is the extension cell ξ

$$\begin{array}{ccc}
 A & \xrightarrow{f_! \otimes f^*} & A \\
 f \downarrow & \rho & \downarrow f \\
 B & \xrightarrow{y_B} & B
 \end{array}
 \qquad
 \begin{array}{ccc}
 A & \xrightarrow{y_A} & A \\
 f \downarrow & \xi & \downarrow f \\
 B & \xrightarrow{f^* \otimes f_!} & B
 \end{array}$$

A morphism $e: A \rightarrow E$ in an equipment is a **cover** if the canonical globular cell is an iso $e^* \otimes e_! \cong y_E$. Dually, a morphism $m: E \rightarrow B$ is an **inclusion** if the canonical globular cell is an iso $m_! \otimes m^* \cong y_E$.

A Factorization System

Theorem

Suppose that \mathbb{D} is a unit-pure cartesian equipment with strong and \mathcal{M} -monic tabulators. If local products satisfy Frobenius and every relation is a tabulator, then with $\mathcal{E} = \text{covers}$ and $\mathcal{M} = \text{inclusions}$, $\mathcal{F} = (\mathcal{E}, \mathcal{M})$ is a proper, stable factorization system on \mathbb{D}_0 .

Given $f: A \rightarrow B$, a factorization arises as in the diagram:

$$\begin{array}{ccc}
 A & \xrightarrow{y} & B \\
 e \downarrow & \exists! & \downarrow e \\
 T(f^* \otimes f_!) & \xrightarrow{y} & T(f^* \otimes f_!) \\
 l \downarrow & \tau & \downarrow r \\
 B & \xrightarrow{f^* \otimes f_!} & B
 \end{array}
 \quad = \quad
 \begin{array}{ccc}
 A & \xrightarrow{y} & B \\
 f \downarrow & \xi & \downarrow f \\
 B & \xrightarrow{f^* \otimes f_!} & B
 \end{array}$$

using the universal property of the tabulator.

Theorem

If \mathbb{D} is a unit-pure, cartesian equipment such that

1. tabulators exist and are strong, \mathcal{M} -monic and functorial,
2. every relation is a tabulator, and finally
3. local products satisfy Frobenius

then the identity functor $1: \mathbb{D}_0 \rightarrow \mathbb{D}_0$ extends to an adjoint equivalence

$$\mathbb{R}el(\mathbb{D}_0; \mathcal{F}) \simeq \mathbb{D}$$

where \mathcal{F} is the orthogonal factorization system given by inclusions and covers.

Theorem

A double category \mathbb{D} is equivalent to one $\mathbb{R}el(\mathcal{C}; \mathcal{F})$ for some proper and stable factorization system on a finitely-complete category \mathcal{C} if, and only if, \mathbb{D} is

1. a unit-pure cartesian equipment;
2. with strong, \mathcal{M} -monic and functorial tabulators;
3. in which every relation is the tabulator of its cokernel;
4. and for which local products satisfy Frobenius.

For more see [Lam21].

THANK YOU!



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